

Applications of Number Theory in Enhancing Biometric Data Privacy through Hash Function

Satriadhikara Panji Yudhistira - 13522125¹

Program Studi Teknik Informatika

Sekolah Teknik Elektro dan Informatika

Institut Teknologi Bandung, Jl. Ganesha 10 Bandung 40132, Indonesia

¹13522125@std.stei.itb.ac.id

Abstract— In this modern world, our unique identities such as fingerprints, iris patterns, and even facial contours are increasingly being used to identify and verify who we are. However, storing this sensitive information in plain text can pose a fatal security threat.

This paper delves into the world of number theory, exploring how its tools - modular arithmetic and - can be leveraged to design efficient hash functions specifically for biometric data protection in computing technology.

Keywords— Biometric Data Privacy, Hash Function, Number Theory.

I. INTRODUCTION

Do you unlock your phone using face? Or fingerprint? That is some examples of authentication using biometric. Why would we use biometric? What if we want to implement it as well to our local authentication like our home door or something like that.

Biometrics (ancient Greek: bios ="life", metron ="measure") refers to two very different fields of study and application. The first, which is the older and is used in biological studies, including forestry, is the collection, synthesis, analysis and management of quantitative data on biological communities such as forests. Biometrics in reference to biological sciences has been studied and applied for several generations and is somewhat simply viewed as "biological statistics" [1].

Our biometric data is one of the most important identification data. It is really needed that we secure our private data (the biometric data) in a safe way. With the application of number theory in hash function, we can securely save or store our biometric data.

II. THEORETICAL BASIS

A. Modular Arithmetic

Modular arithmetic plays an important role in integer computation, especially in application of cryptography. Operator that used in modulo arithmetic is **mod**. Mod operator gives the remainder of a division. For example, 23 divided by 5 is 4 and remainder of 3, so we write $23 \bmod 5 = 3$. The definition of mod operator is stated as follows:

Notation: $a \bmod m = r$ so that $a = mq + r$, with $0 \leq r < m$

If $a \bmod m = 0$, then we can say that a is multiple of m , that is a is divisible by m [1].

B. Hash Function

Data that saved on the computer memory need to place in a way that can be search fast. Every data in the form of record has a unique key field that differentiates from a record and another record. Hash function is used to place a record that have a key. Hash function is commonly used such as

$$h(k) = k \bmod m$$

that in this case, m is a sum of memory location that available (for example if the memory is such as cell that given index of 0 until $m - 1$). The h function above place record with k key into certain memory that address of $h(k)$.

Suppose that $m = 11$, so we have cell of memory that given index of 0 to 10. We will store the record data that every each one has a key 15, 558, 32, 132, 102, and 5. In the beginning the cell of memory is empty.

Each these six data stored on the location that calculated as follows:

$$h(15) = 15 \bmod 11 = 4$$

$$h(558) = 558 \bmod 11 = 8$$

$$h(32) = 32 \bmod 11 = 10$$

$$h(132) = 132 \bmod 11 = 0$$

$$h(102) = 102 \bmod 11 = 3$$

$$h(5) = 5 \bmod 11 = 5$$

The state of the cells of memory after been stored the sixth record data is visualized as follows:

132			102	15	5			558		32
0	1	2	3	4	5	6	7	8	9	10

Because hash function is not one-to-one function, so there may be a collision in storing some record data. For example, we will store a record data with the key is 257. The calculated hash generates $h(257) = 257 \bmod 11 = 4$, whereas the cell memory index of 4 is already filled. So, we got a collision. To overcome a collision, we need to apply collision resolution policy. One collision resolution policy is finding the next highest cell that has not been filled. Then if we apply this policy, the data with the key 257 is stored at cell index of 6.

If we want to find the certain record data, then we use the hash function again. For example, we will find the record data with the key p , then we calculate $h(p) = p \bmod 11$, for instance $h(p) = q$. If the record p is the same as the data in the location of q , we say record p location is found. Otherwise, if the record p is not the same as the data in the location of q , so we look at

REFERENCES

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ATTACHMENT

Source code used in this research:
https://github.com/satriadhikara/biometric_hash

PERNYATAAN

Dengan ini saya menyatakan bahwa makalah yang saya tulis ini adalah tulisan saya sendiri, bukan saduran, atau terjemahan dari makalah orang lain, dan bukan plagiasi.

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Satriadhikara Panji Yudhistira 13522125